

Europäisches Patentamt

European Patent Office

Office européen des brevets



EP 0 730 356 A2

(12)

EUROPEAN PATENT APPLICATION

(43) Date of publication:

04.09.1996 Bulletin 1996/36

(51) Int. Cl.⁶: **H04L 1/00**

(11)

(21) Application number: 96108491.0

(22) Date of filing: 19.01.1993

(84) Designated Contracting States:

AT BE CH DE DK ES FR GB GR IE IT LI LU MC NL PT SE

(30) Priority: 16.01.1992 US 822164

(62) Application number of the earlier application in accordance with Art. 76 EPC: 93903541.6

(71) Applicant: QUALCOMM INCORPORATED San Diego, California 92121 (US)

(72) Inventors:

 Padovani, Roberto San Diego, CA 92130 (US)

- Tiedemann, Edward, G., Jr. San Diego, CA 92122 (US)
- Weaver, Lindsay, A., Jr. Boulder, CO 80303 (US)
- Butler, Brian, K.
 Cardiff, CA 92007 (US)
- (74) Representative: O'Sullivan, Gearoid Patrick et al Tomkins & Co.,5 Dartmouth Road Dublin 6 (IE)

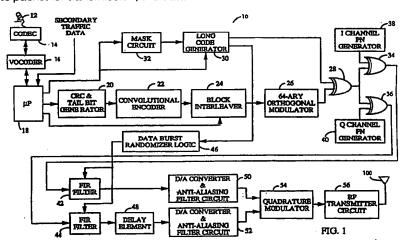
Remarks:

This application was filed on 28 - 05 - 1996 as a divisional application to the application mentioned under INID code 62.

(54) Communication system and method for transmitting data of different types

(57) The invention relates to a communication system in which transmission takes place according to a format which permits different types of data to be combined and transmitted within a single transmission. The novel feature is that the communication system transmits variable length frames of data in packets, and that a data combining and transmission sub-system (14, 16, 18, 20) is provided so that when a frame of data does not require a complete packet for transmission, the data

combining sub-system combines the frame of data with additional data to provide a complete packet. The data combining sub-system comprises input means for receiving the frame of data and the additional data and for combining the frame of data and the additional data to provide a complete packet responsive to a control signal, and control means for providing the control signal.



Description

10

BACKGROUND OF THE INVENTION

I. Field of the Invention

The present application relates to the organization of data for transmission. More particularly, the present invention relates to a novel and improved method and apparatus for formatting vocoder data, non-vocoder data and signaling data for transmission.

II. Description of the Related Art

In the field of digital communications various arrangements of digital data for transmission are used. The data bits are organized according to commonly used formats for transfer over the communication medium.

It is therefore an object of the present invention to provide a data format which facilitates the communication of various types of data, and data of various rates, to be communicated in a structured form.

US-A-4,291,409 discloses a method and apparatus employing spread spectrum techniques in a wide bandwidth communications system. A plurality of transmitting stations are each equipped to provide a transmission signal representing a pseudo-random coded, phase modulated, message signal. The transmission signal is directed through a bandwidth which encompasses otherwise dedicated, relatively narrow bandwidth repeater channels, employed in connection with a communications satellite, to a generally fixed receiver station. At the receiving station, the incoming signal is (a) code acquired and tracked, (b) carrier acquired and tracked, (c) phase locked to the receiver local oscillator and (d) coherently demodulated to extract the desired data. The receiving station advantageously employs plural receiving elements each having a pseudo-random sequence code matched filter which significantly reduces code acquisition time by obviating the necessity of exhaustively correlating the incoming signal with a replica of the pseudo-random code word at the receiver station.

EP-A2-0 412 583 discloses a time division multiple access (TDMA) communication device controller, which controls all signaling, synchronization and supervisory functions. In one embodiment, the invention operates to control a remote communication device having a vocoder and buffering means. The remote communication device is enabled to operate as a dispatch, full duplex or a combination dispatch/full duplex communication device. In another embodiment, a primary station (repeater) is controlled to operate as a single frequency repeater (SFR) or as a multi-frequency TDMA repeater. A communication channel is divided into time sub-slots which may be allocated to different users. The number of sub-slots allocated to a given user depends on the quality of speech required by the user. At the beginning of a call the user selects the required quality of speech. For the duration of the call the speech is encoded at that fixed rate. The encoded speech is provided to a buffer. The buffer is capable of holding enough encoded data to fill a time sub-slot. As soon as the buffer is filled the data is entered into the sub-slot for transmission. If the user wants to provide speech and digital data, this need is specified prior to transmission. The communication system then provides enough sub-slots to carry the fixed rate speech data and also provides additional sub-slots to carry the fixed rate digital data. When the call is completed the caller is billed in accordance with the number of sub-slots used.

WO 91/07030 discloses a distributed synchronization method for a wireless fast packet communication system. The distributed synchronization method, according to the invention, provides for the combination of both voice and data in a single switch using a common packet structure. It allows for the dynamic synchronization of packets. This includes not only bandwidth within the voice or data areas of the frame, but also between the voice and data portions.

The method includes generating a set of tail bits for being appended to data in a frame.

SUMMARY OF THE INVENTION

The present invention is a novel and improved method and system for formatting digital data for communication over a transmission medium.

In communication systems it is important to utilize a data format which permits a full communication of data between users. In a communication system, such as a code division multiple access (CDMA) communication system, in which it is desirable to communicate various types of data, and at various rates, a data format must be selected which permits maximum flexibility within a predefined structure. Furthermore to maximize resources it is desirable to permit a sharing of the format to permit different types of data to be organized together. In such situations it is necessary to structure the data in a manner in which it may be readily extracted according to the corresponding type and rate.

A method and apparatus is provided for arranging various types of data, and at various rates, into a uniquely structured format for transmission. Data is provided as vocoder data or different types of non-vocoder data. The data is organized into frames of a predetermined time duration for transmission. The data frames are organized, depending on the data, to be at one of several data rates. Vocoder data is provided at one of several data rates and is organized in

45

the frame according to a predetermined format. Frames may be formatted with a sharing of vocoder data with non-vocoder data to be at a highest frame data rate. Non-vocoder data may be organized so as to also be at a highest frame rate. Additional control data may be provided within the data frames to support various aspects of the transmission and recovery upon reception.

The invention in its widest aspect is set forth in Claims 1 and 9 written with regard to EP-A2-0 412 583 as nearest prior art.]

BRIEF DESCRIPTION OF THE DRAWINGS

5

10

15

20

25

30

35

The features, objects, and advantages of the present invention will become more apparent from the detailed description set forth below when taken in conjunction with the drawings in which like reference characters identify correspondingly throughout and wherein:

Figure 1 is a block diagram illustrating an exemplary embodiment for a transmitter portion of a transceiver;

Figures 2a - 2h are a series of diagrams illustrating frame data formats for the various data rates, types and modes;

Figure 3 is a diagram illustrating an exemplary circuit implementation of the CRC and Tail Bit generator of Figure 1;

Figures 4a - 4e is a flow chart of the formatting of frames of data;

Figures 5a - 5d illustrate in a series of charts the ordering of code symbols in the interleaver array for transmission data rates of 9.6, 4.8, 2.4 and 1.2 kbps, respectively;

Figures 6a - 6c is a chart illustrating the Walsh symbol corresponding to each encoder symbol group;

Figure 7 is a block diagram illustrating the long code generator of Figure 1;

Figures 8a - 8c are a series of diagrams illustrating long code masks for the various channel type; and

Figure 9 is a graph illustrating the frequency response of the digital filters of Figure 1.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

Referring now to the drawings, Figure 1 illustrates an exemplary embodiment of a transmit portion 10 of a CDMA mobile station transceiver or PCN handset. In a CDMA cellular communication system a forward CDMA channel is used to transmit information from a cell base station to the mobile station. Conversely a reverse CDMA channel is used to transmit information from the mobile station to the cell base station. The communication of signals from the mobile station may be characterized in the form of an access channel or a traffic channel communication. The access channel is used for short signaling messages such as call originations, responses to pages, and registrations. The traffic channel is used to communicate (1) primary traffic, typically includes user speech, or (2) secondary traffic, typically user data, or (3) signaling traffic, such as command and control signals, or (4) a combination of primary traffic and secondary traffic or (5) a combination of primary traffic and signaling traffic.

Transmit portion 10 enables data to be transmitted on the reverse CDMA channel at data rates of 9.6 kbps, 4.8 kbps, 2.4 kbps or 1.2 kbps. Transmissions on the reverse traffic channel may be at any of these data rates while transmissions on the access channel are at the 4.8 kbps data rate. The transmission duty cycle on the reverse traffic channel will vary with the transmission data rate. Specifically, the transmission duty cycle for each rate is provided in Table I. As the duty cycle for transmission varies proportionately with the data rate, the actual burst transmission rate is fixed at 28,800 code symbols per second. Since six code symbols are modulated as one of 64 Walsh symbols for transmission, the Walsh symbol transmission rate shall be fixed at 4800 Walsh symbols per second which results in a fixed Walsh chip rate of 307.2 kcps.

All data that is transmitted on the reverse CDMA channel is convolutional encoded, block interleaved, modulated by 64-ary modulation, and direct-sequence PN spread prior to transmission. Table I further defines the relationships and rates for data and symbols for the various transmission rates on the reverse traffic channel. The numerology is identical for the access channel except that the transmission rate is fixed at 4.8 kbps, and the duty cycle is 100%. As described later herein each bit transmitted on the reverse CDMA channel is convolutional encoded using a rate 1/3 code. Therefore, the code symbol rate is always three times the data rate. The rate of the direct-sequence spreading functions shall be fixed at 1.2288 MHz, so that each Walsh chip is spread by precisely four PN chips.

TABLE I

Bit Rate (kbps)	9.6	4.8	2.4	1.2
PN Chip Rate (Mcps)	1.2288	1.2288	1.2288	1.2288
Code Rate (bits/code symbol)	1/3	1/3	1/3	1/3
TX Duty Cycle (%)	100.0	50.0	25.0	12.5
Code Symbol Rate (sps)	28,800	28,800	28,800	28,800
Modulation (code symbol/Walsh symbol)	6	6	6	6
Walsh Symbol Rate (sps)	4800	4800	4800	4800
Walsh Chip; Rate (kcps)	307.20	307.20	307.20	307.20
Walsh Symbol (μs)	208.33	208.33	208.33	208.33
PN Chips/Code Symbol	42.67	42.67	42.67	42.67
PN Chips/Walsh Symbol	256	256	256	256
PN Chips/Walsh Chip	4	4	4	4

Transmit portion 10, when functioning in mode in which primary traffic is present, communicates acoustical signals, such as speech and/or background noise, as digital signals over the transmission medium. To facilitate the digital communication of acoustical signals, these signals are sampled and digitized by well known techniques. For example, in Figure 1, sound is converted by microphone 12 to an analog signal which is then converted to a digital signal by codec 14. Codec 14 typically performs an analog to digital conversion process using a standard 8 bit/µlaw format. In the alternative, the analog signal may be directly converted to digital form in a uniform pulse code modulation (PCM) format. In an exemplary embodiment codec 14 uses an 8 kHz sampling and provides an output of 8 bit samples at the sampling rate so as to realize a 64 kbps data rate.

The 8-bit samples are output from codec 14 to vocoder 16 where a µlaw/uniform code conversion process is performed. In vocoder 16, the samples are organized into frames of input data wherein each frame is comprised of a predetermined number of samples. In a preferred implementation of vocoder 16 each frame is comprised of 160 samples or of 20 msec. of speech at the 8 kHz sampling rate. It should be understood that other sampling rates and frame sizes may be used. Each frame of speech samples is variable rate encoded by vocoder 16 with the resultant parameter data formatted into a corresponding data packet. The vocoder data packets are then output to microprocessor 18 and associated circuitry for transmission formatting. Microprocessor 18 generically includes program instructions contained with a program instruction memory, a data memory, and appropriate interface and related circuitry as is known in the art.

A preferred implementation of vocoder 16 utilizes a form of the Code Excited Linear Predictive (CELP) coding techniques so as to provide a variable rate in coded speech data. A Linear Predictive Coder (LPC) analysis is performed upon a constant number of samples, and the pitch and codebook searches are performed on varying numbers of samples depending upon the transmission rate. A variable rate vocoder of this type is described in further detail in WO 92/22891.

Vocoder 16 may be implemented in an application specific integrated circuit (ASIC) or in a digital signal processor. In the variable rate vocoder just mentioned, the speech analysis frames are 20 msec. in length, implying that the extracted parameters are output to microprocessor 18 in a burst 50 times per second. Furthermore the rate of data output is varied from roughly 8 kbps to 4 kbps to 2 kbps, and to 1 kbps.

At full rate, also referred to as rate 1, data transmission between the vocoder and the microprocessor is at an 8.55 kbps rate. For the full rate data the parameters are encoded for each frame and represented by 160 bits. The full rate data frame also includes a parity check of 11 bits thus resulting in a full rate frame being comprised of a total of 171 bits. In the full rate data frame, the transmission rate between the vocoder and the microprocessor absent the parity check bits would be 8 kbps.

At half rate, also referred to as rate 1/2, data transmission between the vocoder and the microprocessor is at a 4 kbps rate with the parameters encoded for each frame using 80 bits. At quarter rate, also referred to as rate 1/4, data transmission between the vocoder and the microprocessor is at a 2 kbps rate with the parameters encoded for each frame using 40 bits. At eighth rate, also referred to as rate 1/8, data transmission between the vocoder and the microprocessor is slightly less than a 1 kbps rate with the parameters encoded for each frame using 16 bits.

10

15

20

40

In addition, no information may be sent in a frame between the vocoder and the microprocessor. This frame type, referred to as a blank frame, may be used for signaling or other non-vocoder data.

The vocoder data packets are then output to microprocessor 18 and CRC and Tail Bit generator 20 for completing the transmission formatting. Microprocessor 18 receives packets of parameter data every 20 msec. along with a rate indication for the rate the frame of speech samples was encoded. Microprocessor 18 also receives, if present, an input of secondary traffic data for output to generator 20. Microprocessor 18 also internally generates signaling data for output to generator 20. Data whether it is primary traffic, secondary traffic or signaling traffic matter, if present, is output from microprocessor 18 to generator 20 every 20 msec. frame.

Generator 20 generates and appends at the end of all full and half rate frames a set of parity check bits or cyclic redundancy check bits (CRC Bits) which are used at the receiver as a frame quality indicator. For a full rate frame, regardless of whether the data is a full rate primary, secondary or signaling traffic, or a combination of half rate primary and secondary traffic, or a combination of half rate primary and signaling traffic, generator 20 preferably generates a set of CRC Bits according to a first polynomial. For a half rate data frame, generator 20 also generates a set of CRC Bits preferably according to a second polynomial. Generator 20 further generates for all frame rates a set of Encoder Tail Bits which follow the CRC bits, if present, or data if CRC bits are not present, at the end of the frame. Further details of the operation on microprocessor 18 and generator 20 are provided later herein with reference to Figures 3 and 4.

Reverse traffic channel frames provided from generator 20 at the 9.6 kbps rate are 192 bits in length and span the 20 msec. frame. These frames consist of a single Mixed Mode Bit, auxiliary format bits if present, message bits, a 12-bit frame quality indicator (CRC), and 8 Encoder Tail Bits as shown in Figures 2a - 2e. The Mixed Mode Bit shall be set to '0' during any frame in which the message bits are primary traffic information only. When the Mixed Mode Bit is '0', the frame shall consist of the Mixed Mode Bit, 171 Primary Traffic bits, 12 CRC Bits, and 8 Encoder Tail Bits.

The Mixed Mode Bit is set to '1' for frames containing secondary or signaling traffic. In these instances the first bit following the Mixed Mode Bit is a Burst Format Bit which specifies whether the frame is in a "blank-and-burst" or a "dimand-burst" format. A "blank-and-burst" operation is one in which the entire frame is used for secondary or signaling traffic while a "dim-and-burst" operation is one in which the primary traffic shares the frame with either secondary or signaling traffic. If the Burst Format Bit is a '0', the frame is of the "dim and burst format", and if a '1' the frame is of the "blank and burst format".

The second bit following the Mixed Mode Bit is a Traffic Type Bit. The Traffic Type Bit is used to specify whether the frame contains secondary or signaling traffic. If the Traffic Type Bit is a '0', the frame contains signaling traffic, and if a '1', the frame contains secondary traffic. Figures 2b - through 2e illustrate the Burst Format Bit and the Traffic Type Bit.

When the Burst Format Bit is '0' denoting dim-and-burst, the two bits following the Traffic Type Bit are Traffic Mode Bits. These bits indicate the number of bits that are used for primary traffic information and the number of bits that shall be used for either signaling or secondary traffic information within that frame. For a default mode, only the Traffic Mode '00' is defined with all other traffic modes reserved for other bit type and numbers. Referring to Figures 2b and 2c, in the exemplary and preferred embodiment, 80 bits are used for primary traffic (half rate vocoder data packet) while 86 and 87 bits are respectively used for signaling and secondary traffic.

In frames where there is signaling traffic present the first bit of the frame's signaling portion is a Start of Message (SOM) Bit. The SOM Bit is a '1' if a reverse traffic channel message (signaling message) begins at the following bit. Generally the first bit of a reverse traffic channel message does not begin anywhere else in the frame other than following the SOM Bit. However should the frame contain part of a message that began in a previous frame the SOM Bit is a '0'. If the SOM Bit is a '0' the following bit is part of the message but it is not the first bit of the complete message.

In the preferred implementation only primary traffic is transmitted in frames at the 4.8 kbps, 2.4 kbps, and 1.2 kbps rates. Mixed mode operation is generally not be supported at rates other than the 9.6 kbps rate, although it may be readily configured to do so. The frame formats for these particular rates are shown in Figures 2f - 2h. For the 4.8 kbps rate, the frame is 96 bits in length with the bits spaced over the 20 msec. time period of the frame as described later herein. The 4.8 kbps rate frame contains 80 primary traffic bits, an 8-bit frame quality indicator (CRC), and 8 Encoder Tail Bits. For the 2.4 kbps rate, the frame is 48 bits in length with the bits spaced over the 20 msec. time period of the frame as also described later herein. The 2.4 kbps rate frame contains 40 primary traffic bits and 8 Encoder Tail Bits. For the 1.2 kbps rate, the frame is 24 bits in length with the bits spaced over the 20 msec. time period of the frame as also described later herein. The 1.2 kbps rate frame contains 16 primary traffic bits and 8 encoder tail bits.

In a preferred embodiment the access channel data is generated by microprocessor 18 for transmission at a rate of 4.8 kbps. As such the data is prepared in a manner identical to that of 4.8 kbps frame format data, such as encoding, interleaving as Walsh encoding. In the encoding scheme implemented for the 4.8 kbps data, whether reverse traffic channel data or access channel data, redundant data is generated. Unlike the reverse traffic channel where the redundant data is eliminated in the transmission, in the access channel all data including redundant data is transmitted. Details on the transmission aspects of frames of access channel data are provided later herein.

Figure 3 illustrates an exemplary implementation of the elements for formatting the data in accordance with Figures 2a - 2h. In Figure 3 data is transmitted from microprocessor 18 (Figure 1) to generator 20. Generator 20 is comprised of data buffer and control logic 60, CRC circuits 62 and 64, and Tail Bit circuit 66. Along with data provided from the

microprocessor a rate command may optionally be provided. Data is transferred for each 20 msec frame from the microprocessor to logic 60 where temporarily stored. For each frame, logic 60 may for each frame count the number of bits transmitted from the microprocessor, or in the alternative use the rate command and a count of the clock cycles in formatting a frame of data.

Each frame of the traffic channel includes a frame quality indicator. For the 9.6 kbps and 4.8 kbps transmission rates, the frame quality indicator is the CRC. For the 2.4 kbps and 1.2 kbps transmission rates, the frame quality indicator is implied, in that no extra frame quality bits are transmitted. The frame quality indicator supports two functions at the receiver. The first function is to determine the transmission rate of the frame, while the second function is to determine whether the frame is in error. At the receiver these determinations are made by a combination of the decoder information and the CRC checks.

For the 9.6 kbps and 4.8 kbps rates, the frame quality indicator (CRC) is calculated on all bits within the frame, except the frame quality indicator (CRC) itself and the Encoder Tail Bits. Logic 60 provides the 9.6 kbps and 4.8 kbps rate data respectively to CRC circuits 62 and 64. Circuits 62 and 64 are typically constructed as a sequence of shift registers, modulo-2 adders (typically exclusive-OR gates) and switches as illustrated.

The 9.6 kbps transmission rate data uses a 12-bit frame quality indicator (CRC), which is be transmitted within the 192-bit long frame as discussed with reference to Figures 2a - 2e. As illustrated in Figure 3 for CRC circuit 62, the generator polynomial for the 9.6 kbps rate is as follows:

$$g(x) = x^{12} + x^{11} + x^{10} + x^{9} + x^{8} + x^{4} + x + 1.$$
 (1)

The 4.8 kbps transmission rate data uses an 8-bit CRC, which is transmitted within the 96-bit long frame as discussed with reference to Figure 2f. As illustrated in Figure 3 for CRC circuit 64, the generator polynomial for the 4.8 kbps rate is as follows:

$$g(x) = x^{8} + x^{7} + x^{4} + x^{3} + x + 1.$$
 (2)

Initially, all shift register elements of circuits 62 and 64 are set to logical one ('1') by an initialization signal from logic 60. Furthermore logic 60 set the switches of circuits 62 and 64 in the up position.

For 9.6 kbps rate data, the registers of circuit 62 are then clocked 172 times for the 172 bits in the sequence of primary traffic, secondary traffic or signaling bits or a mixture thereof along with the corresponding mode/format indicator bits as input to circuit 62. After 172 bits are clocked through circuit 62, logic 60 then sets the switches of circuit 62 in the down position with the registers of circuit 62 then being clocked an additional 12 times. As a result of the 12 additional clockings of circuit 62, 12 additional output bits are generated which are the CRC bits. The CRC bits, in the order calculated, are appended to the end of the 172 bits as output from circuit 62. It should be noted that the 172 bits output from logic 60 which pass through circuit 62 are undisturbed by the computation of the CRC bits and are thus output from circuit 62 in the same order and at the same value at which they entered.

For 9.6 kbps rate data bits are input to circuit 64 from logic 60 in the following order. For the case of primary traffic only, the bits are input to circuit 64 from logic 60 in the order of the single mixed mode (MM) bit followed by the 171 primary traffic bits. For the case of "dim and burst" with primary and signaling traffic, the bits are input to circuit 64 from logic 60 in the order of the single MM bit, a single burst format (BF) bit, a traffic type (TT) bit, a pair of traffic mode (TM) bits, 80 primary traffic bits, a start of message (SOM) bit, and 86 signaling traffic bits. For the case of "dim and burst" with primary and secondary traffic, the bits are input to circuit 64 from logic 60 in the order of the single MM bit, the single BF bit, the TT bit, the pair of TM bits, 80 primary traffic bits and 87 signaling traffic bits. For the case of "blank and burst" data format with signaling traffic only, the bits are input to circuit 64 from logic 60 in the order of the single MM bit, the single BF bit, the TT bit, the SOM bit and 168 signaling traffic bits. For the case of "blank and burst" data format with secondary traffic only, the bits are input to circuit 64 from logic 60 in the order of the single MM bit, the single BF bit, the TT bit and 169 signaling traffic bits.

Similarly for 4.8 kbps rate data, the registers of circuit 64 are clocked 80 times for the 80 bits of primary traffic data, or for the 80 bits of access channel data, as input to circuit 64 from logic 60. After the 80 bits are clocked through circuit 64, logic 60 then sets the switches of circuit 64 in the down position with the registers of circuit 64 then being clocked an additional 8 times. As a result of the 12 additional clockings of circuit 62, 12 additional output bits are generated which are the CRC bits. The CRC bits, in the order calculated, are again appended to the end of the 80 bits as output from circuit 64. It should again be noted that the 80 bits output from logic 60 which pass through circuit 64 are undisturbed by the computation of the CRC bits and are thus output from circuit 64 in the same order and at the same value at which they entered.

The bits output from either of circuits 62 and 64 are provided to switch 66 which is under the control of logic 60. Also input to switch 66 are the 40 and 16 bits of primary traffic data output from logic 60 for 2.4 kbps and 1.2 kbps data frames. Switch 66 selects between providing an output of the input data (up position) and tail bits at a logical zero ('0') value (down position). Switch 66 is normally set in the up position to permit data from logic 60, and from circuits 62 and

15

20

64 if present, to be output from generator 20 to encoder 22 (Figure 1). For the 9.6 kbps and 4.8 kbps frame data, after the CRC bits are docked through switch 66, logic 60 sets the switch to the down position for 8 clock cycles so as to generate 8 all zero tail bits. Thus for 9.6 kbps and 4.8 kbps data frames, the data as output to the encoder for the frame includes appended after the CRC bits, the 8 tail bits. Similarly for the 2.4 kbps and 1.2 kbps frame data, after the primary traffic bits are clocked from logic 60 through switch 66, logic 60 sets the switch to the down position for 8 clock cycles so as to again generate 8 all zero tail bits. Thus for 2.4 kbps and 1.2 kbps data frames, the data as output to the encoder for the frame includes appended after the primary traffic bits, the 8 tail bits.

Figures 4a - 4e illustrate in a series of flow charts the operation of microprocessor 18, and generator 20 in assembling the data into the disclosed frame format. It should be noted that various schemes may be implemented for giving the various traffic types and rates priority for transmission. In an exemplary implementation, when a signaling traffic message is to be sent when there is vocoder data present a "dim and burst" format may be selected. Microprocessor 18 may generate a command to vocoder 18 for the vocoder to encode speech sample frames at the half rate, regardless of the rate at which the vocoder would normally encode the sample frame. Microprocessor 18 then assembles the half rate vocoder data with the signaling traffic into the 9.6 kbps frame as illustrated in Figure 2b. In this case, a limit may be place on the number of speech frames encoded at the half rate to avoid degradation in the speech quality. In the alternative, microprocessor 18 may wait until a half rate frame of vocoder data is received before assembling the data into the "dim and burst" format. In this case, in order to ensure timely transmission of the signaling data, a maximum limit on the number of consecutive frames at other than half rate may be imposed before a command is sent to the vocoder to encode at half rate. Secondary traffic may be transferred in the "dim and burst" format (Figure 2c) in a similar manner.

Similar is the case for the "blank and burst" data formats as illustrated in Figures 2d - 2d. The vocoder may be commanded to not encode the frame of speech samples or the vocoder data is ignored by the microprocessor in constructing the data frame. Prioritizing between generating frame formats of primary traffic of various rate, "dim and burst" traffic, and "blank and burst" traffic is open to many possibilities.

Referring back to Figure 1, 20 msec. frames of 9.6 kbps, 4.8 kbps, 2.4 kbps and 1.2 kbps data are thus output from generator 20 to encoder 22. In the exemplary embodiment encoder 22 is a preferably a convolutional encoder, a type of encoder well known in the art. Encoder 22 preferably encodes the data using a rate 1/3, constraint length k = 9 convolutional code. As an example encoder 22 is constructed with generator functions of $g_0 = 557$ (octal), $g_1 = 663$ (octal) and $g_2 = 711$ (octal). As is well known in the art, convolutional encoding involves the modulo-2 addition of selected taps of a serially time-shifted delayed data sequence. The length of the data sequence delay is equal to k-1, where k is the code constraint length. Since in the preferred embodiment a rate 1/3 code is used, three code symbols, the code symbols (c_0) , (c_1) and (c_2) , are generated for each data bit input to the encoder. The code symbols (c_0) , (c_1) and (c_2) are respectively generated by the generator functions g_0 , g_1 and g_2 . The code symbols are output from encoder 22 to block interleaver 24. The output code symbols are provided to interleaver 24 in the order of the code symbol (c_0) being first, the code symbol (c_1) being second and the code symbol (c_2) being last. The state of the encoder 22, upon initialization, is the all-zero state. Furthermore the use of tail bits at the end of each frame provides a resetting of encoder 22 to an all-zero state.

The symbols output from encoder 22 are provided to block interleaver 24 which under the control of microprocessor 18 provides a code symbol repetition. Using a conventional random access memory (RAM) with the symbols stored therein as addressed by microprocessor 18, code symbols may be stored in a manner to achieve a code symbol repetition rate that varies with the data channel.

Code symbols are not be repeated for the 9.6 kbps data rate. Each code symbol at the 4.8 kbps data rate is repeated 1 time, i.e. each symbol occurs 2 times. Each code symbol at the 2.4 kbps data rate is repeated 3 times, i.e. each symbol occurs 4 times. Each code symbol at the 1.2 kbps data rate is repeated 7 times, i.e. each symbol occurs 8 times. For all data rates (9.6, 4.8, 2.4 and 1.2 kbps), the code repetition results in a constant code symbol rate of 28,800 code symbols per second for the data as output from interleaver 24. On the reverse traffic channel the repeated code symbols are not transmitted multiple times with all but one of the code symbol repetitions deleted prior to actual transmission due to the variable transmission duty cycle as discussed in further detail below. It should be understood that the use of code symbol repetition as an expedient method for describing the operation of the interleaver and a data burst randomizer as discussed again in further detail below. It should be further understood that implementations other than those that use code symbol repetition may be readily devised that achieve the same result and remain within the teaching of the present invention.

All code symbols to be transmitted on the reverse traffic channel and the access channel are interleaved prior to modulation and transmission. Block interleaver 24, constructed as is well known in the art, provides an output of the code symbols over a time period spanning 20 msec. The interleaver structure is typically a rectangular array with 32 rows and 18 columns, i.e. 576 cells. Code symbols are written into the interleaver by columns, with repetition for data at the 9.6, 4.8, 2.4 and 1.2 kbps rate, so as to completely fill the 32 × 18 matrix. Figures 5a - 5d illustrate the ordering of write operations of repeated code symbols into the interleaver array for transmission data rates of 9.6, 4.8, 2.4 and 1.2 kbps, respectively.

20

Reverse traffic channel code symbols are output from the interleaver by rows. Microprocessor 18 also controls the addressing of the interleaver memory for outputting the symbols in the appropriate order. The interleaver rows are preferably output in the following order:

5 At 9.6 kbps:

10

20

40

50

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32

At 4.8 kbps:

1 3 2 4 5 7 6 8 9 11 10 12 13 15 14 16 17 19 18 20 21 23 22 24 25 27 26 28 29 31 30 32

At 2.4 kbps:

; 1 5 2 6 3 7 4 8 9 13 10 14 11 15 12 16 17 21 18 22 19 23 20 24 25 29 26 30 27 31 28 32

At 1.2 kbps:

1 9 2 10 3 11 4 12 5 13 6 14 7 15 8 16 17 25 18 26 19 27 20 28 21 29 22 30 23 31 24 32.

Access channel code symbols are also output from interleaver 24 by rows. Microprocessor 18 again controls the addressing of the interleaver memory for outputting the symbols in the appropriate order. The interleaver rows are output in the following order at the 4.8 kbps rate for the access channel code symbols:

1 17 9 25 5 21 13 29 3 19 11 27 7 23 15 31 2 18 10 26 6 22 14 30 4 20 12 28 8 24 16 32.

It should be noted that other encoding rates, such as a rate 1/2 convolutional code used on the forward transmission channel, along with various other symbol interleaving formats may be readily devised using the basic teaching of the present invention

Referring again to Figure 1, the interleaved code symbols are output from interleaver 24 to modulator 26. In the preferred embodiment modulation for the Reverse CDMA Channel uses 64-ary orthogonal signaling. That is, one of 64 possible modulation symbols is transmitted for each six code symbols. The 64-ary modulation symbol is one of 64 orthogonal waveforms generated preferably using Walsh functions. These modulation symbols are given in Figures 6a - 6c and are numbered 0 through 63 The modulation symbols are selected according to the following formula:

Modulation symbol number =
$$c_0 + 2_{c1} + 4_{c2} + 8_{c3} + 16_{c4} + 32_{c5}$$
 (3)

where c_5 shall represent the last or most recent and c_0 the first or oldest binary valued ('0' and '1') code symbol of each group of six code symbols that form a modulation symbol. The period of time required to transmit a single modulation symbol is referred to as a "Walsh symbol" interval and is approximately equal to 208.333 μ s. The period of time associated with one-sixty-fourth of the modulation symbol is referred to as a "Walsh chip" and is approximately equal to 3.2552083333... μ s.

Each modulation or Walsh symbol is output from modulator 26 to one input of a modulo-2 adder, exclusive-OR gate 28. The Walsh symbols are output from modulator at a 4800 sps rate which corresponds to a Walsh chip rate of 307.2 kcps. The other input to gate 28 is provided from long code generator 30 which generates a masked pseudonoise (PN) code, referred to as the long code sequence, in cooperation with mask circuit 32. The long code sequence provided from generator 30 is at a chip rate rate four times the Walsh chip rate of modulator 26, i.e. a PN chip rate 1.2288 Mcps. Gate 28 combines the two input signals to provide an output of data at the chip rate of 1.2288 Mcps.

The long code sequence is a time shift of a sequence of length 2⁴²-1 chips and is generated by a linear generator well known in the art using the following polynomial:

$$p(x) = x^{42} + x^{35} + x^{33} + x^{31} + x^{27} + x^{26} + x^{25} + x^{22} + x^{21} + x^{19}$$

$$+ x^{18} + x^{17} + x^{16} + x^{10} + x^{7} + x^{6} + x^{5} + x^{3} + x^{2} + x^{1} + 1$$
(4)

Figure 7 illustrates generator 30 in further detail. Generator 30 is comprised of a sequence generator section 70 and a masking section 72. Section 70 is comprised of a sequence of shift registers and modulo-2 adders (typically exclusive-OR gates) coupled together to generate a 42-bit code according to equation 4. The long code is then generated by masking the 42-bit state variables output from section 70 with a 42-bit wide mask provided from mask circuit 32.

Section 72 is comprised of a series of input AND gates $74_1 - 74_{42}$ having one input for receiving a respective mask bit of the 42-bit wide mask. The other input of each of AND gates $74_1 - 74_{42}$ receives the output from a corresponding

shift register in section 70. The output of AND gates $74_1 - 74_{42}$ are modulo-2 added by adder 76 to form a single bit output for each 1.2288 MHz clocking of the shift registers of section 70. Adder 76 is typically constructed as a cascaded arrangement of exclusive-OR gates as is well known in the art. Therefore, the actual output PN sequence is generated by the modulo-2 addition of all 42 masked output bits of sequence generator 70 as shown in Figure 7.

The mask used for the PN spreading shall vary depending on the channel type on which the mobile station is communicating. Referring to Figure 1, an intialization information is provided from microprocessor 18 to generator 30 and circuit 32. Generator 30 is responsive to the initialization information for initialization of the circuitry. Mask 32 is also responsive to the initialization information, which indicates the mask type to be provided, to output a 42-bit mask. As such, mask circuit 32 may be configured as a memory which contains a mask for each communication channel type. Figures 8a - 8c provide an exemplary definition of the masking bits for each channel type.

Specifically, when communicating on the Access Channel, the mask is defined as illustrated in Figure 8a. In the Access Channel mask, mask bits M_{24} through M_{41} are set to '1'; mask bits M_{19} through M_{23} are set to the chosen Access Channel number; mask bits M_{16} through M_{18} are set to the code channel for the associated Paging Channel, i.e, the range typically being 1 through 7; mask bits M_{9} through M_{15} are set to the registration zone; for the current base station; and mask bits M_{0} through M_{8} are set to the pilot PN value for the current CDMA Channel.

When communicating on the Reverse Traffic Channel, the mask is defined as illustrated in Figure 8b. The mobile station uses one of two long codes unique to that mobile station: a public long code unique to the mobile station's electronic serial number (ESN); and a private long code unique for each mobile identification number (MIN) which is typically the telephone number of the mobile station. In the public long code the mask bits M_{32} through M_{41} are set to '0,' and the mask bits M_0 through M_{31} are set to the mobile station ESN value.

It is further envisioned that a private long code may be implemented as illustrated in Figure 8c. The private long code will provide additional security in that it will only be known to the base station and the mobile station. The private long code will not be transmitted in the clear over the transmission medium. In the private long code the mask bit M_{40} through M_{41} are set to '0' and '1' respectively; while mask bits M_0 through M_{39} may be set to according to a predetermined assignment scheme.

Referring back to Figure 1 the output of gate 28 is respectively provided as one input to each one of a pair of modulo-2 adders, exclusive-OR gates 34 and 36. The other input to each of gates 34 and 36 are second and third PN sequences are I and Q channel "short codes" respectively generated by I and Q Channel PN generators 38 and 40. The Reverse Access Channel and Reverse Traffic Channel is therefore OQPSK spread prior to actual transmission. This offset quadrature spreading on the Reverse Channel uses the same I and Q PN codes as the Forward Channel I and Q pilot PN codes. The I and Q PN codes generated by generators 38 and 40 are of length 2¹⁵ and are preferably the zero-time offset codes with respect to the Forward Channel. For purposes of further understanding, on the Forward Channel a pilot signal is generated for each base station. Each base station pilot channel signal is spread by the I and Q PN codes as just mentioned. Base station I and Q PN codes are offset from one another, by a shifting of the code sequence, so as to provide distinction between base station transmission. The generating functions for the I and Q short PN codes shall be as follows:

$$P_1(x) = x^{15} + x^{13} + x^9 + x^8 + x^7 + x^5 + 1$$
 (5)

o and

$$P_Q(x) = x^{15} + x^{12} + x^{11} + x^{10} + x^6 + x^5 + x^4 + x^3 + 1.$$
 (6)

Generators 38 and 40 may be constructed as is well known in the art so as to provide an output sequence in accordance with equations (5) and (6).

The I and Q waveforms are respectively output from gates 34 and 36 where respectively provided as inputs to finite impulse response (FIR) filters 42 and 44. FIR filters 42 and 44 are digital filters which bandlimit the resulting I and Q waveforms. These digital filters shape the I and Q waveforms such that the resulting spectrum is contained within a given spectral mask. The digital filters preferably have the impulse response shown in the following Table II:

TABLE II

h(0) =	-0.02204953170628	= h(46)		h(12) =	0.03881898337058	= h(34)
h(1) =	-0.01997721494122	= h(45)		h(13) =	0.10411392223653	= h(33)
h(2) =	-0.00905191683798	= h(44)]	h(14) =	0.11268193747141	= h(32)
h(3) =	0.02005789896688	= h(43)]	h(15) =	0.04184165339577	= h(31)
h(4) =	0.05926358628876	= h(42)]	h(16) =	-0.08271278252498	= h(30)
h(5) =	0.09021366056377	= h(41)		h(17) =	-0.18998156787345	= h(29)
h(6) =	0.09304356333555	= h(40)		h(18) =	-0.19486048259840	= h(28)
h(7) =	0.05917668051274	= h(39)		h(19) =	-0.04343248005925	= h(27)
h(8) =	0.00032251394639	= h(38)		h(20) =	0.25121616493295	= h(26)
h(9) =	-0.05381152911745	= h(37)		h(21) =	0.60403450701992	= h(25)
h(10) =	-0.07036222587323	= h(36)		h(22) =	0.89017616954958	= h(24)
h(11) =	-0.03405975708422	= h(35)		h(23) =	1	= h(23)

Filters 42 and 44 may be constructed according to well known digital filter techniques and preferably provide a frequency response as illustrated in Figure 9.

The binary '0' and '1' inputs to digital filters 42 and 44, generated by the PN spreading functions, are mapped into +1 and -1, respectively. The sampling frequency of the digital filter is 4.9152 MHz = 4×1.2288 MHz. An additional binary '0' and '1' input sequence synchronous with the I and Q digital waveforms shall be provided to each of digital filters 42 and 44. This particular sequence, referred to as a masking sequence, is the output generated by a data burst randomizer. The masking sequence multiplies the I and Q binary waveforms to produce a ternary (-1, 0, and +1) input to the digital filters 42 and 44.

As discussed previously the data rate for transmission on the Reverse Traffic Channel is at one of the rates of equal 9.6, 4.8, 2.4, or 1.2 kbps and varies on a frame-by-frame basis. Since the frames are of a fixed 20 ms length for both the Access Channel and the Reverse Traffic Channel, the number of information bits per frame shall be 192, 96, 48, or 24 for transmission at data rates of 9.6, 4.8, 2.4, or 1.2 kbps, respectively. As described previously, the information is encoded using a rate 1/3 convolutional encoder and then the code symbols shall be repeated by a factor of 1, 2, 4, or 8 for a data rate of 9.6, 4.8, 2.4, or 1.2 kbps, respectively. The resulting repetition code symbol rate is thus fixed at 28,800 symbols per second (sps). This 28,800 sps stream is block interleaved as previously described.

Prior to transmission, the Reverse Traffic Channel interleaver output stream is gated with a time filter that allows transmission of certain interleaver output symbols and deletion of others. The duty cycle of the transmission gate thus varies with the transmit data rate. When the transmit data rate is 9.6 kbps, the transmission gate allows all interleaver output symbols to be transmitted. When the transmit data rate is 4.8 kbps, the transmission gate allows one-half of the interleaver output symbols to be transmitted, and so forth. The gating process operates by dividing the 20 msec frame into 16 equal length (i.e., 1.25 ms) periods, called power control groups. Certain power control groups are gated on (i.e., transmitted), while other groups are gated off (i.e., not transmitted).

The assignment of gated-on and gated-off groups is referred to as a data burst randomizer function. The gated-on power control groups are pseudo-randomized in their positions within the frame so that the actual traffic load on the Reverse CDMA Channel is averaged, assuming a random distribution of the frames for each duty cycle. The gated-on power control groups are such that every code symbol input to the repetition process shall be transmitted once without repetition. During the gated-off periods, the mobile station does not transmit energy, thus reducing the interference to other mobile stations operating on the same Reverse CDMA Channel. This symbol gating occurs prior to transmission filtering.

The transmission gating process is not used when the mobile station transmits on the Access Channel. When transmitting on the Access Channel, the code symbols are repeated once (each symbol occurs twice) prior to transmission.

In the implementation of the data burst randomizer function, data burst randomizer logic 46 generates a masking stream of 0's and 1's that randomly mask out the redundant data generated by the code repetition. The masking stream pattern is determined by the frame data rate and by a block of 14 bits taken from the long code sequence generated by generator 30. These mask bits are synchronized with the data flow and the data is selectively masked by these bits through the operation of the digital filters 42 and 44. Within logic 46 the 1.2288 MHz long code sequence output from

5

10

15

20

25

45

generator 30 is input to a 14-bit shift register, which is shifted at a 1.2288 MHz rate. The contents of this shift register are loaded into a 14-bit latch exactly one power control group (1.25 ms) before each Reverse Traffic Channel frame boundary. Logic 46 uses this data along with the rate input from microprocessor 18, to determine, according to a predetermined algorithm, the particular power control group(s) in which the data is to be allowed to pass through filters 42 and 46 for transmission. Logic 46 thus outputs for each power control group a '1' or '0' for the entire power control group depending on whether the data is to be filtered out ('0') or passed through ('1'). At the corresponding receiver, which also uses the same long code sequence and a corresponding rate determined for the frame, determines the appropriate power control group(s) in which the data is present.

The I channel data output from filter 42 is provided directly to a digital to analog (D/A) converter and anti-aliasing filter circuit 50. The Q channel data however is output from filter 44 to a delay element 48 which a one-half PN chip time delay (406.9 nsec) in the Q channel data. The Q channel data is output from delay element 48 to digital to analog (D/A) converter and anti-aliasing filter circuit 52. Circuits 50 and 52 convert the digital data to analog form and filter the analog signal. The signals output from circuits 50 and 52 are provided to Offset Quadrature Phase Shift Key (OQPSK) modulator 54 where modulated and output to RF transmitter circuit 56. Circuit 56 amplifies, filters and frequency upconverts the signal for transmission. The signal is output from circuitry 56 to antenna 58 for communication to the base station via transmitter 100.

It should be understood that the exemplary embodiment of the present invention discusses the formatting of data for modulation and transission with respect to a mobile station. It should be understood that the data formatting is the same for a cell base station, however the modulation may be different.

Claims

20

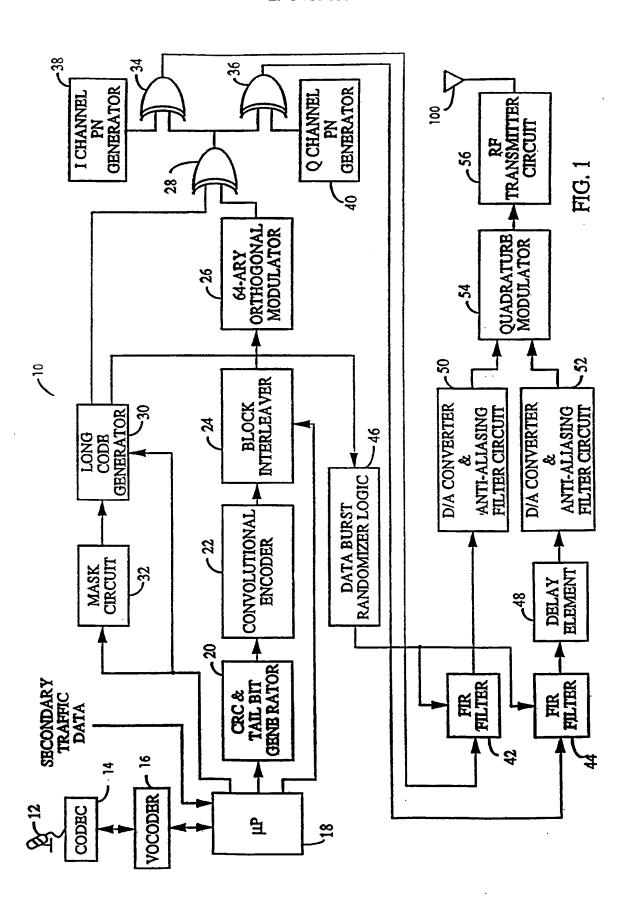
25

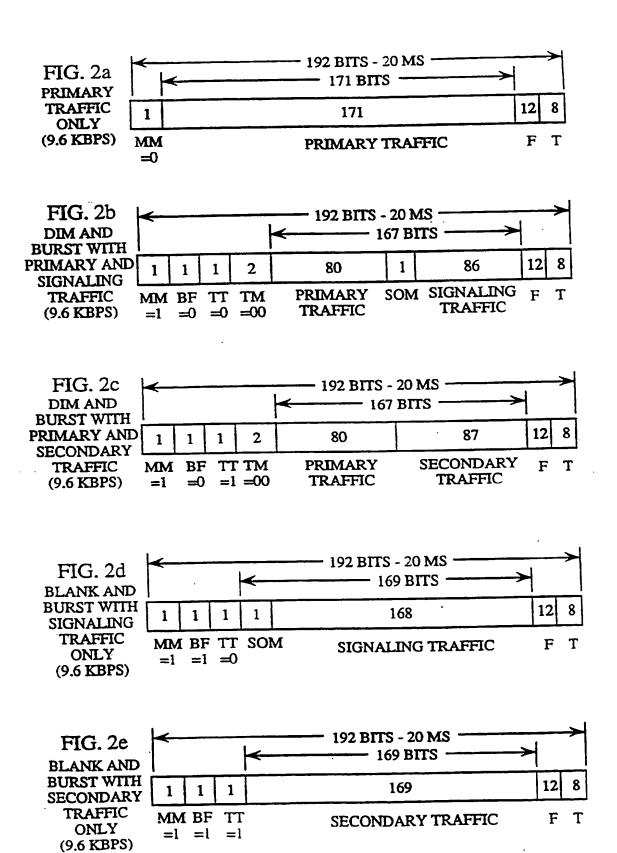
30

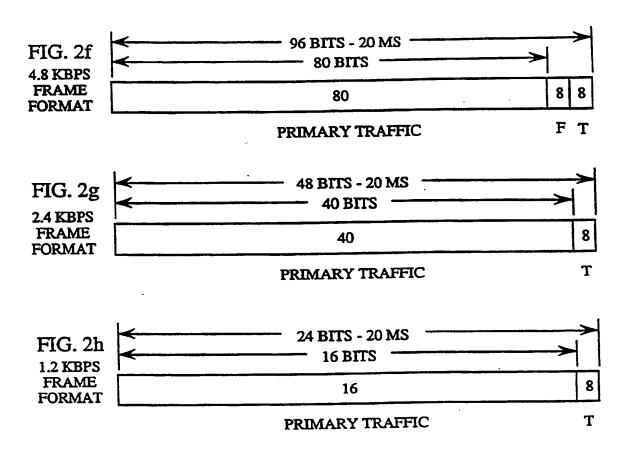
- 1. A communication system in which transmission takes place according to a format which permits different types of data to be combined and transmitted within a single transmission, characterized in that the communication system transmits variable length frames of data in packets, and in that a data combining and transmission sub-system (14, 16, 18, 20) is provided so that when a frame of data of said variable length frames of data does not require a complete packet for transmission, the data combining sub-system combines said frame of data with additional data to provide a complete packet, said data combining sub-system comprising input means for receiving said frame of data and additional data and for combining said frame of data and said additional data to provide said complete packet responsive to a control signal, and control means for providing said control signal.
- 2. A communication system according to Claim 1 wherein said control means (18) is responsive to a data rate signal.
- 3. A communication system according to Claim 1 wherein said frame of data comprises speech data and said additional data comprises signalling data.
 - 4. A communication system according to Claim 1 wherein said frame of data comprises speech data and said additional data comprises secondary traffic data.
- 40 5. A communication system according to any one of Claims 1 to 4 wherein said transmission packet further comprises at least one overhead bit (TT) indicative of a type of said additional data.
 - A communication system according to any of Claims 1 to 5 which is a digital communication system.
- 45 7. A communication system according to any of Claims 1 to 5 which is a spread spectrum communications system.
 - 8. A communication system according to any of the preceding claims in which the input means and control means of said sub-system comprise:
 - variable rate vocoder means (16) for receiving samples of speech data, encoding said speech samples to provide coded speech data at a data rate of a plurality of data rates;
 - processor means (18) for receiving said coded speech data and additional data, and when said speech data is provided at a rate less than a predetermined maximum, combining said coded speech data with said additional data to provide a packet of data.
- 9. A method for use in a communication system for transmitting data of a first type and data of a second type in packets having a data capacity including the steps of receiving said data of said first type and said data of said second type characterized in that when said data of said first type does not use all of said data capacity in a data packet, the excess capacity in said data packet is used by combining data of said second type with said data of said first type to provide said packet.

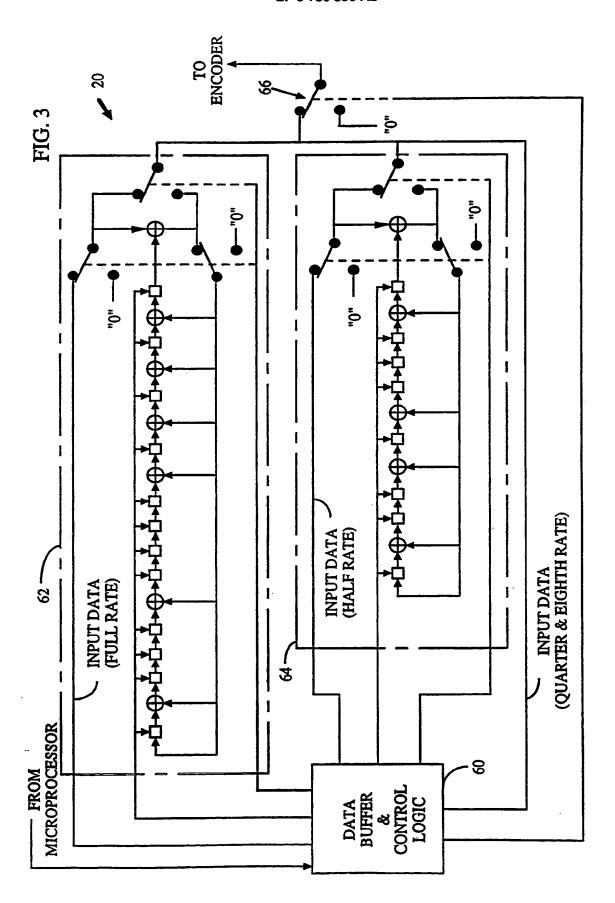


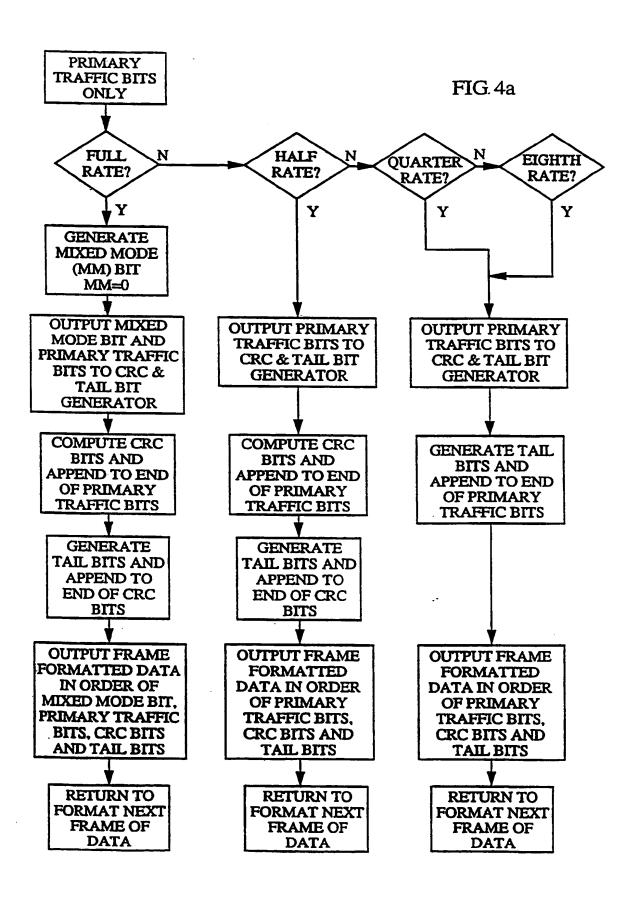
- 10. The method of Claim 9 wherein said data of said first type comprises coded speech data and said data of said second type comprises signaling data.
- 11. The method of Claim 10 wherein said data of said first type comprises coded speech data and said data of said second type comprises secondary traffic data.
 - 12. The method of any of Claims 9 to 11 when used in a digital communication system.
- 13. The method of any of Claims 9 to 11 when used in a spread spectrum communications system, the combined data being transmitted in accordance with a spread spectrum modulation format.

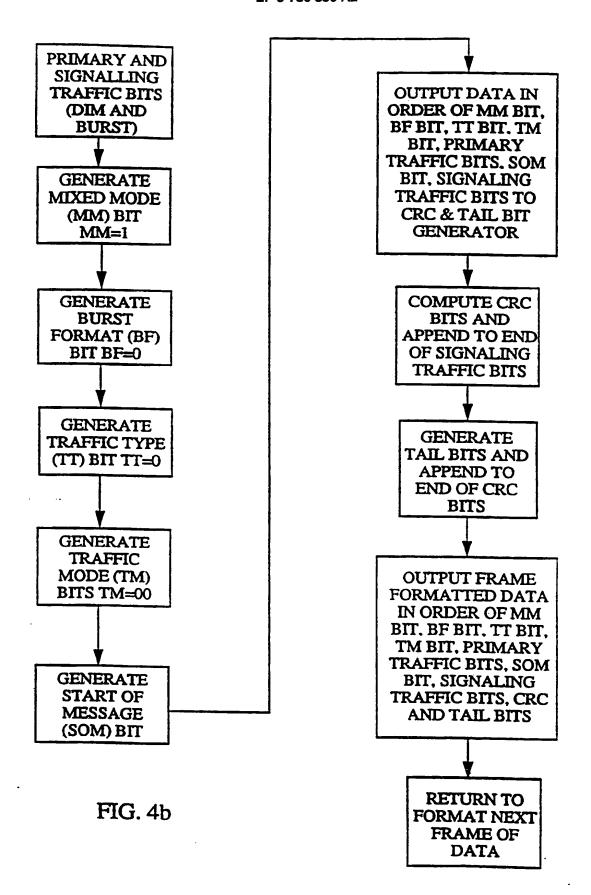


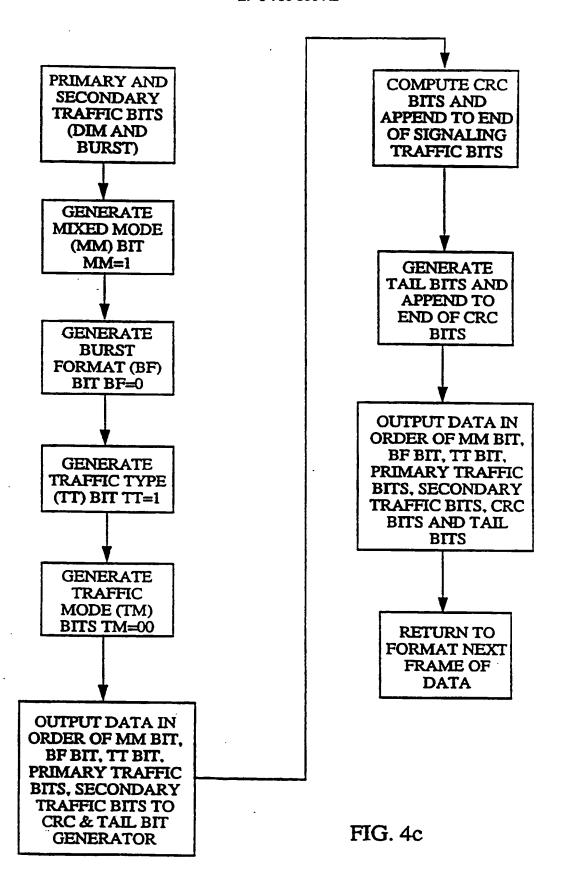


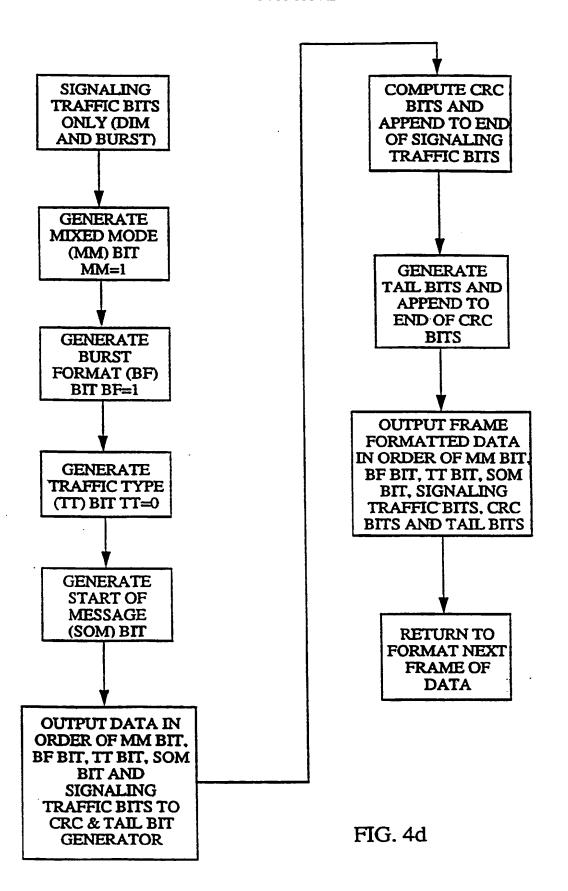












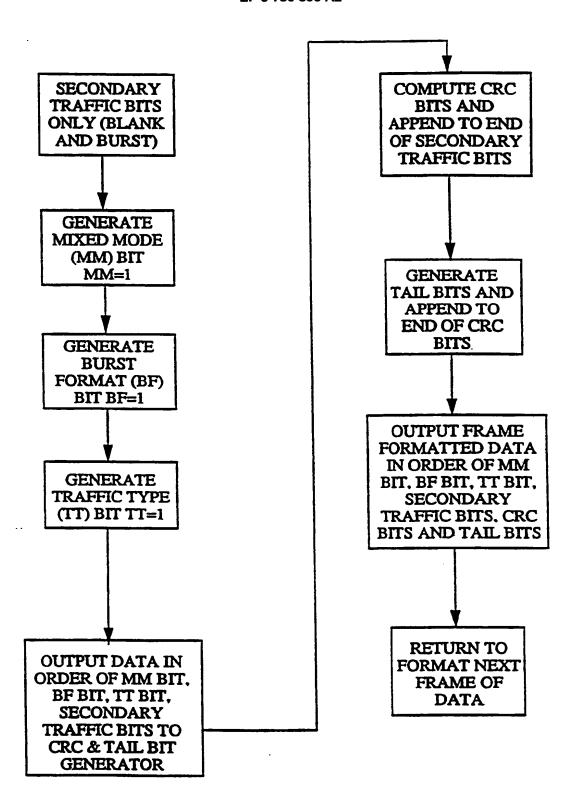


FIG. 4e

65 97 129 161 193 225 257 289 321 353 385 417 449 481 513 545 66 98 130 162 194 226 258 290 322 354 386 418 450 482 514 546 67 99 131 163 195 227 259 291 323 355 387 419 451 483 515 547 68 100 132 164 196 228 260 292 324 356 388 420 452 484 516 548 69 101 133 165 197 229 261 293 325 357 389 421 453 485 517 549 70 102 134 166 198 230 262 294 326 358 390 422 454 486 518 550 36 37 6 38 263 295 264 296 265 297 266 298 327 359 328 360 329 361 330 362 7 39 71 103 135 167 199 231 391 423 455 487 519 551 72 104 136 168 200 232 40 392 424 456 488 520 9 41 73 105 137 169 201 233 393 425 394 426 457 489 10 74 106 138 170 202 234 42 458 490 75 107 139 171 203 235 267 299 76 108 140 172 204 236 268 300 77 109 141 173 205 237 269 301 11 43 331 363 364 395 427 459 491 44 12 332 333 396 428 460 492 524 556 13 45 365 397 429 461 493 525 557 206 238 270 302 334 366 398 430 462 494 526 558 207 239 271 303 335 367 399 431 463 495 527 559 208 240 272 304 336 368 400 432 464 496 528 560 14 46 78 110 142 174 15 16 79 111 143 175 47 48 80 112 144 176 17 209 241 273 305 337 369 401 433 465 497 529 561 210 242 274 306 338 370 402 434 466 498 530 562 49 81 113 145 177 18 50 82 114 146 178 83 115 147 179 211 243 275 307 339 371 403 435 467 499 531 563 84 116 148 180 212 244 276 308 340 372 404 436 468 500 532 564 85 117 149 181 213 245 277 309 341 373 405 437 469 501 533 565 19 51 20 52 21 53 214 246 278 310 342 374 406 438 470 502 534 566 215 247 279 311 343 375 407 439 471 503 535 567 216 248 280 312 344 376 408 440 472 504 536 568 217 249 281 313 345 377 409 441 473 505 537 569 218 250 282 314 346 378 410 442 474 506 538 570 219 251 283 315 347 379 411 443 475 507 539 571 220 252 284 316 348 380 412 444 476 508 540 572 22 23 24 54 55 56 57 58 59 86 118 150 182 87 119 151 183 88 120 152 184 89 121 153 185 90 122 154 186 91 123 155 187 25 26 27 28 92 60 124 156 188 29 93 125 157 189 61 221 253 285 317 381 413 445 477 509 541 573 349 94 126 158 190 222 254 286 318 350 382 414 446 478 510 542 574 95 127 159 191 223 255 287 319 351 383 415 447 479 511 543 575 96 128 160 192 224 256 288 320 352 384 416 448 480 512 544 576 30 62 63

FIG. 5a

209 225 241 257 113 129 145 161 177 193 34 161 177 193 162 178 194 209 225 210 226 50 66 82 113 129 145 114 130 146 35 51 210 226 211 227 83 98 114 130 146 162 178 194 99 115 131 147 163 179 195 67 67 68 68 69 242 258 243 259 36 36 37 37 38 52 52 53 53 54 99 115 131 147 163 179 195 211 227 243 259 20 84 100 116 132 148 100 116 132 148 164 180 196 164 180 196 212 228 212 228 244 260 244 260 213 229 165 181 197 101 117 133 149 245 261 165 181 197 213 229 245 166 182 198 214 230 246 22 22 23 70 86 101 117 133 149 102 118 134 150 39 55 71 102 118 134 150 103 119 135 151 166 182 198 167 183 199 214 230 246 215 231 247 87 103 119 135 151 167 183 199 215 231 247 263 168 184 200 168 184 200 216 232 216 232 24 40 72 104 120 136 152 104 120 136 152 248 56 88 217 233 249 169 185 105 121 137 217 233 217 233 218 234 218 234 219 235 219 235 26 74 154 169 185 201 170 186 202 250 58 90 105 121 106 122 138 170 186 202 170 186 202 171 187 203 171 187 203 172 188 204 172 188 204 75 75 251 27 91 138 154 139 155 283 43 59 106 122 107 123 267 107 123 139 155 28 220 236 220 236 252 284 12 44 76 92 108 124 140 156 108 124 140 156 268 221 237 141 157 173 189 109 125 30 94 173 189 174 190 206 221 237 222 238 46 78 109 125 110 126 141 157 142 158 62 15 222 238 223 239 255 271 31 79 95 142 158 143 159 175 191 207 47 63 127 223 239 111 127 143 159 175 191 256 176 192 208 176 192 208 224 240 224 240 64 112 128 144 160 144 160

FIG. 5b

73 33 33 34 34 34 35 35 35 36 36 37 37 37 37 38 38 25 26 26 26 27 27 27 27 28 28 29 29 29 29 30 97 105 113 121 129 137 97 105 113 121 129 137 97 105 113 121 129 137 74 74 74 58 58 66 66 90 90 90 90 91 91 97 105 113 121 129 137 50 50 82 98 106 114 122 130 138 42 43 43 43 44 44 44 45 45 45 98 106 114 122 130 138 98 106 114 122 130 138 98 106 114 122 130 138 99 107 115 123 131 139 99 107 115 123 131 139 83 83 51 51 59 59 75 75 76 76 76 77 77 77 77 78 99 107 115 123 131 139 91 99 107 115 123 131 139 91 99 107 115 123 131 139 92 100 108 116 124 132 140 92 100 108 116 124 132 140 92 100 108 116 124 132 140 92 100 108 116 124 132 140 52 52 52 53 53 53 54 54 55 55 55 55 55 60 60 60 84 84 84 84 68 68 68 20 20 20 20 21 21 12 101 109 117 125 133 141 69 70 70 101 109 117 125 133 141 85 21 22 22 22 23 23 23 24 101 109 117 125 133 141 93 101 109 117 125 133 141 94 102 110 118 126 134 142 94 102 110 118 126 134 142 70 78 79 79 79 63 63 63 94 102 110 118 126 134 142 94 102 110 118 126 134 142 39 39 95 103 111 119 127 135 143 71 71 72 72 95 103 111 119 127 135 143 95 103 111 119 127 135 143 95 103 111 119 127 135 143 104 112 120 128 136 144 56 96 104 112 120 128 136 144 96 104 112 120 128 136 144 96 104 112 120 128 136 144

FIG. 5c

21 21 21 21 21 21 25 25 25 25 26 26 26 26 27 27 27 27 27 27 28 28 28 28 28 28 29 29 29 29 29 30 30 30 3333333444444445535535353566 53 53 53 53 53 53 54 54 54 13 13 13 57 57 57 57 57 57 58 58 58 17 13 13 37 38 22 22 22 22 22 22 23 23 23 23 23 23 24 46 50 50 50 50 50 38 38 38 42 18 42 42 42 42 54 54 54 54 18 18 18 19 62 58 58 59 59 59 6. 38 38 50 50 51 51 51 30 31 31 31 31 31 31 32 67 67 67 47 39 39 39 43 43 43 55 19 19 19 19 20 20 20 63 15 59 59 59 59 55 55 55 39 43 51 51 51 15 15 63 12 52 52 52 32 32 32 32 24 24 24 24 20 20 20 44 52 72

FIG. 5d

FIG. 6a

WALSH CHIP WITHIN SYMBOL

								WAL	H CH	T W	SIL	360	اد				
L	Γ			=	1111	1111	2222	2222	2233	3333	3333	4444	444	4455	5555	5555	9999
		0123	4567	8901	2345	6489	0123	4567	8901	2345	6289	0123	4567	8901	2345	6789	0123
	c		000	0000	0000	0000	800	9000 9000	0000	0000	0000	0000	0000	0000	8	8	000
	·-	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101	0101
	. ~	9	00	0011	0011	0011	0011	0011	0011	0011	8011	8011	9011	8011	8011	8011	9
	٠,	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110	0110
	7	3	1111	000	1111	000	1111	000	1111	000	1111	<u>000</u>	1111	8	1111	8	1111
	- ~	200	1010	010	1010	0101	1010	0101	1010	0101	1010	0101	1010	0101	1010	0101	1010
	2	35	100	0	1100	00	1100	0011	1100	0011	1100	911	1100	0011	1100	0011	1100
	,	0110		0110	1001	0110	1001	0110	1001	0110	1001	0110	1001	0110	1001	0110	1001
	- 0			1111	1111	000	000	1111	1111	0000	000	1111	1111	800	8	1111	1111
_	0	300		1010	1010	0101	0101	1010	1010	0101	0101	1010	1010	0101	0101	1010	1010
-	\ <u> </u>	3 5		120	1100	50	0011	1100	1100	0011	0011	1100	1100	0011	9011	1100	1100
_	2 =		0110		1001	0110	0110	1001	1001	0110	0110	1001	1001	0110	0110	1001	1001
	2:		1111	1111	000	0000	1111	1111	000	000	1111	1111	8 8	800	1111	1111	8
	3 5	200	1010	1010	0101	0101	1010	1010	0101	0101	1010	1010	0101	0101	1010	1010	0101
	7	011	222	100	5	9	1100	100	0011	0011	1100	1188	<u>8</u>	<u>8</u>	1188	1100	0011
	7	0110	1001	19	0110	0110	1001	1001	0110	0110	1001	1001	0110	0110	1001	<u>8</u>	0110
	2 2					1111	1111	1111	1111	000	8	8	8	1111	1111	1111	1111
	2 5	35		010	010	1010	1010	1010	1010	0101	0101	0101	0101	1010	1010	1010	1010
	<u> </u>	35		1100	00	100	1100	1100	1100	0011	0011	0011	0011	118	1188	118	118
	2 0	121	100	0110	0110	1001	100	1001	1001	0110	0110	0110	0110	1001	1001	1001	1001
	3 :	35	1111		1111	1111	000	1111	000	8	1111	000	1111	1111	8	1111	88
,	; ;		1010	010	1010	1010	0101	1010	0101	0101	1010	0101	1010	1010	0101	1010	010
	3 :	5	100	00	1100	1100	0011	1100	0011	0011	100	001	1100	1100	0011	118	00
	3 1	0110	196	0110	1001	1001	0110	1001	0110	0110	1001	0110	1001	1001	0110	<u> </u>	0110
ك																	

MAINE SHEDY -ZUEX

V)
C)
日	

11 1111 1111 2222 2222 2233 3333 444 4444 4455 5555 5555 5666 24 6000 6000 1111 1111 1111 1111 0000 0000 0000 0000 1111 111			•					WAL	SHCHI	IP WIT	HINS	YMBO	L				
0123 4567 8901 2345 6789 0123 4567 8901 2345 6789 0123 4567 8901 2345 6789 0000 0000 1111 1111 1111 1111 0000 0000 0000 0000 1111 1111 1111 1111 1111 0000 0101 0101 1010 1010 1010 1010 0101 0101 0101 0101 1100 1100 1100 1100 1100 011 011 0101 1010 1010 1010 1010 0110 0110 0110 0110 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1110 0101 1010 10	Ŀ			F	1111	1111	2222	2222	2233	3333	3333	4444	4444	4455	5555	5555	9999
0000 0000 1111 1000 1010 0101 0101 1010 1100 1100 1100 1100 1100 1100 1100 1100 1101 1010 1101 1010 1101 1010 <th< th=""><th></th><td>0123</td><td>4567</td><td>8901</td><td>2345</td><td>64.89</td><td>0123</td><td>4567</td><td>8901</td><td>2345</td><td>62/9</td><td>0123</td><td>4567</td><td>8901</td><td>2345</td><td>6789</td><td>0123</td></th<>		0123	4567	8901	2345	64.89	0123	4567	8901	2345	62/9	0123	4567	8901	2345	6789	0123
0101 0101 1010 1010 1010 0101 0101 0101 1010 1010 1010 0101 0101 0101 1010 1010 1010 0110 0111 0011 1001 <th< th=""><th>72</th><th>0000</th><th>0000</th><th>1111</th><th>1111</th><th>1111</th><th>H</th><th>0000 0000</th><th>0000 0000</th><th>0000</th><th>0000</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>8</th><th>88</th></th<>	72	0000	0000	1111	1111	1111	H	0000 0000	0000 0000	0000	0000	1111	1111	1111	1111	8	88
0011 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1101 1001 <th< th=""><th>25</th><th>0101</th><th>0101</th><th>1010</th><th>1010</th><th>1010</th><th>1010</th><th>0101</th><th>0101</th><th>0101</th><th>0101</th><th>1010</th><th>1010</th><th>1010</th><th>1010</th><th>0101</th><th>0101</th></th<>	25	0101	0101	1010	1010	1010	1010	0101	0101	0101	0101	1010	1010	1010	1010	0101	0101
0110 0110 1001 1001 1001 1001 0110 0110 0110 0110 1001 1001 1001 1001 1001 0110 0000 0111 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 0111 0000 0111 0010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0111 1110 0110 1100 0111 1100 0111 1100 0111 1100 0111 1100 0111 0110 0	26	0011	00	1100	1100	1100	1100	0011	0011	0011	0011	1100	1100	1100	1100	0011	8011
000 1111 1111 000 1111 000 1111 000 1111 000 1111 000 1111 000 1111 000 1111 000 0101 1010 <th>27</th> <th>0110</th> <th>0110</th> <th>1001</th> <th>1001</th> <th>1001</th> <th>1001</th> <th>0110</th> <th>0110</th> <th>0110</th> <th>0110</th> <th>1001</th> <th>1001</th> <th>1001</th> <th>1001</th> <th>0110</th> <th>0110</th>	27	0110	0110	1001	1001	1001	1001	0110	0110	0110	0110	1001	1001	1001	1001	0110	0110
0110 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 <th< th=""><th>% </th><th>000</th><th>1111</th><th>1111</th><th>0000</th><th>1111</th><th>000</th><th>000</th><th>1111</th><th>8</th><th>1111</th><th>1111</th><th>800</th><th>1111</th><th>8</th><th>000 000</th><th>1111</th></th<>	% 	000	1111	1111	0000	1111	000	000	1111	8	1111	1111	800	1111	8	000 000	1111
0011 1100 1001 1100 1001 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1101 1100 1101 1101 1100 1101 1101 1100 1100 1101 1100 1100 1101 1100 1101 1101 1100 1101 1100 1101 1100 1101 1100 1101 1100 1101 1100 1101 1100 1101 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 <td< th=""><th>29</th><th>010</th><th>1010</th><th>1010</th><th>0101</th><th>1010</th><th>0101</th><th>0101</th><th>1010</th><th>0101</th><th>1010</th><th>1010</th><th>0101</th><th>1010</th><th>0101</th><th>0101</th><th>1010</th></td<>	29	010	1010	1010	0101	1010	0101	0101	1010	0101	1010	1010	0101	1010	0101	0101	1010
0110 1001 1001 1001 1010 <th< th=""><th><u>چ</u></th><th>00</th><th>1100</th><th>100</th><th>00</th><th>1100</th><th>00</th><th>0011</th><th>1100</th><th>0011</th><th>1100</th><th>1100</th><th>0011</th><th>1100</th><th>0011</th><th>0011</th><th>1100</th></th<>	<u>چ</u>	00	1100	100	00	1100	00	0011	1100	0011	1100	1100	0011	1100	0011	0011	1100
0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0010 0101 0101 1010 <th< th=""><th>3</th><th>0110</th><th>1001</th><th>1001</th><th>0011</th><th>1001</th><th>0110</th><th>0110</th><th>1001</th><th>0110</th><th>1001</th><th>1001</th><th>0110</th><th>1001</th><th>0110</th><th>0110</th><th><u>1</u>8</th></th<>	3	0110	1001	1001	0011	1001	0110	0110	1001	0110	1001	1001	0110	1001	0110	0110	<u>1</u> 8
0101 0100 1100 1100 1100 1100 1100 1100 1100 1100 1101 0001 1001 <th< th=""><th>3</th><th>2</th><th></th><th></th><th>000</th><th>0000</th><th>0000</th><th>000</th><th>000</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th></th<>	3	2			000	0000	0000	000	000	1111	1111	1111	1111	1111	1111	1111	1111
0011 0011 0011 0011 0011 0011 0011 001	3 %			010	010	0101	0101	0101	0101	1010	1010	1010	1010	1010	1010	1010	1010
0100 0110 0110 0110 0110 0110 0110 011	3 2	35		55	51	011	8	00	0011	1100	1100	1100	1100	1100	1100	1100	1100
0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0000 1111 0010 1010 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 <th< th=""><th>3,5</th><td></td><td>0110</td><td>0110</td><td>0110</td><td>0110</td><td>0110</td><td>0110</td><td>0110</td><td>1001</td><td>1001</td><td>1001</td><td>1001</td><td>1001</td><td>1001</td><td>1001</td><td>100</td></th<>	3,5		0110	0110	0110	0110	0110	0110	0110	1001	1001	1001	1001	1001	1001	1001	100
0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1010 0101 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1101 1001 0110 1001 1010 <th< th=""><th>3 %</th><td>000</td><td>1111</td><td></td><td>1111</td><td>0000</td><td>1111</td><td>000</td><td>1111</td><td>1111</td><td>800</td><td>1111</td><td>000</td><td>1111</td><td>8</td><td>1111</td><td>88</td></th<>	3 %	000	1111		1111	0000	1111	000	1111	1111	800	1111	000	1111	8	1111	88
0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 0011 1100 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1011 1010 <td< th=""><th>32</th><td></td><td>1010</td><td>010</td><td>1010</td><td>0101</td><td>1010</td><td>0101</td><td>1010</td><td>1010</td><td>0101</td><td>1010</td><td>0101</td><td>1010</td><td>0101</td><td>1010</td><td>0101</td></td<>	32		1010	010	1010	0101	1010	0101	1010	1010	0101	1010	0101	1010	0101	1010	0101
0110 1001 0110 1001 0110 1001 0110 1001 0110 1001 0110 1001 0110 1001 0110 1001 1011 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1110 1100 1100 1110 1110 1110 1110 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 1000 1111 1110 1110 1110 1110 <th< th=""><th>~ ~</th><td>91</td><td>130</td><td>00</td><td>1100</td><td>0011</td><td>1100</td><td>0011</td><td>1100</td><td>1100</td><td>0011</td><td>1100</td><td>0011</td><td>1100</td><td>011</td><td>1100</td><td>9011</td></th<>	~ ~	91	130	00	1100	0011	1100	0011	1100	1100	0011	1100	0011	1100	011	1100	9011
0000 0111 1111 0000 0111 1111 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0000 1111 1111 0010 0101 1010 0101 1010 1010 0101 1010 0110 <th< th=""><th>36</th><td>0110</td><td>1001</td><td>0110</td><td>1001</td><td>0110</td><td>1001</td><td>0110</td><td>1001</td><td>1001</td><td>0110</td><td>1001</td><td>0110</td><td>1001</td><td>0110</td><td>1001</td><td>0110</td></th<>	36	0110	1001	0110	1001	0110	1001	0110	1001	1001	0110	1001	0110	1001	0110	1001	0110
0101 0101 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1010 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1010 0011 1010 1010 1010 1010 0011 1010 1011 1011 1010 1011 1011 1010 1011 <th< th=""><th>8</th><th>000</th><th>000</th><th>1111</th><th>1111</th><th>000</th><th>000</th><th>1111</th><th>1111</th><th>1111</th><th>1111</th><th>8</th><th>8 8 8</th><th>1111</th><th>1111</th><th>8</th><th>88</th></th<>	8	000	000	1111	1111	000	000	1111	1111	1111	1111	8	8 8 8	1111	1111	8	88
0011 0011 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1100 1111 1111 0000 0000 1111 1111 0000 0000 0010 0110 0110 0110 0110 0101 0101 0101 0101 0101 0101 0101 0101 0101 0101 0110 <th< th=""><th>4</th><th>010</th><th>010</th><th>1010</th><th>1010</th><th>0101</th><th>0101</th><th>1010</th><th>1010</th><th>1010</th><th>1010</th><th>0101</th><th>0101</th><th>1010</th><th>1010</th><th>0101</th><th>0101</th></th<>	4	010	010	1010	1010	0101	0101	1010	1010	1010	1010	0101	0101	1010	1010	0101	0101
0110 0110 1001 1001 1001 1001 1001 1010 1010 1010 1010 1010 1010 1010 1010 1011 1111 0000 0000 1011 1111 0000 <th< th=""><th>42</th><th>9</th><th>9</th><th>100</th><th>1100</th><th>0011</th><th>0011</th><th>1100</th><th>1100</th><th>1100</th><th>1100</th><th>0011</th><th>0011</th><th>118</th><th>118</th><th>8011</th><th><u>8</u></th></th<>	42	9	9	100	1100	0011	0011	1100	1100	1100	1100	0011	0011	118	118	8011	<u>8</u>
0000 1111 1111 0000 0000 1111 1111 0000 <th< th=""><th>43</th><th>0110</th><th>0110</th><th>1001</th><th>1001</th><th>0110</th><th>0110</th><th>1001</th><th>1001</th><th>1001</th><th>1001</th><th>0110</th><th>0110</th><th>1001</th><th>1001</th><th>0110</th><th>0110</th></th<>	43	0110	0110	1001	1001	0110	0110	1001	1001	1001	1001	0110	0110	1001	1001	0110	0110
1010 1010 0101 0101 1010 1010 0101 1010 0101 0101 1010 1010 0101 0101 1100 1100 0111 1100 1100 0011 1100 0111 1100 1100 0111 1100 1100 1100 1100 1100 0110 1100 1	4	000	1111	1111	0000	0000	1111	1111	8	1111	8	8	1111	1111	8	8	1111
1100 1100 0011 0011 1100 1100 0011 1100 0011 1100 1100 1100 0011 0011 1001 1001 0110 0110 0110	45	010	1010	1010	0101	0101	1010	1010	0101	1010	0101	0101	1010	1010	0101	0101	1010
1001 1001 0110 0110 1001 1001 0110 1001 0110 0110 1001 1001 0110 0110	46	9	1100	138	00	0011	1100	1100	0011	1100	8011	0011	1100	1100	8011	<u>8</u>	1188
	47	0110	1001	1001	0110	0110	1001	1001	0110	1001	0110	0110	1001	1001	0110	0110	1001

XEDZ- LOBXX HOLV



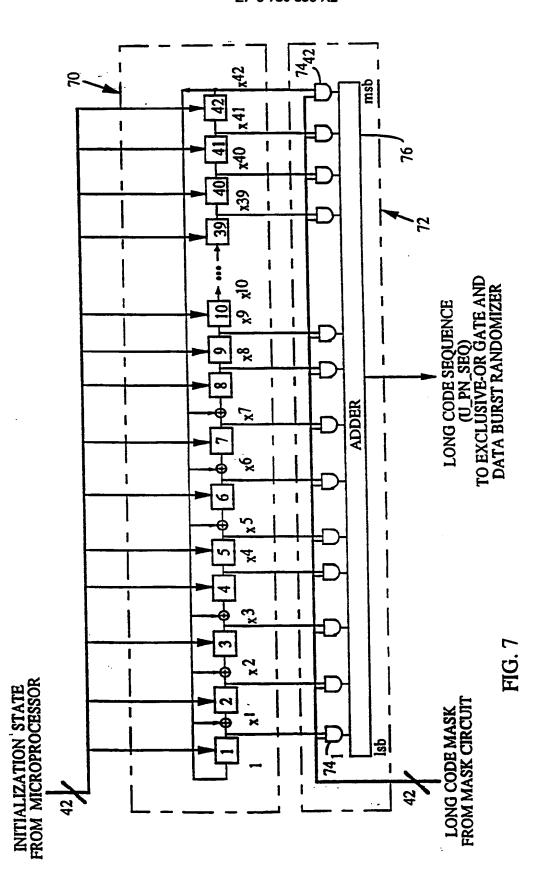
							WALSI	1 CHIP	WITH	N SY	MBOL					
L			II	1111	11111		2222	2233	3333	3333	4444	4444	4455	5555	5555	9999
0	1123	4567	8901	2345	6489	0123	4567	8901	2345	6489	0123	4567	8901	2345	6289	0123
尸	000	0000	0000	0000	1111		IIII	11111	1111	1111	1111	1111	0000	0000	0000	0000
_	101	0101	0101	0101	1010	_	1010	1010	1010	1010	1010	1010	0101	0101	0101	0101
_	110	0011	0011	0011	1100	_	1100	1100	1100	118	1 8	118	811	811	811	8011
_	110	0110	0110	0110	1001		1001	1001	1001	1001	1001	1001	0110	0110	0110	0110
	000	1111	0000	1111	1111		1111	000	1111	0000	=======================================	000	9 9 9	1111	000 000	1111
_	101	1010	0101	1010	1010		1010	0101	1010	0101	1010	0101	0101	1010	0101	1010
_	011	1100	0011	1100	1100		1100	8	1100	011	1188	811	011	118	811	1100
_	110	1001	0110	1001	1001	_	1001	0110	1001	0110	100 100	0110	0110	<u>1</u>	0110	189
	000	0000	1111	1111	1111		000	0000	1111	1111	000	0000	000	000	1111	1111
_	100	0101	1010	1010	1010	_	0101	0101	1010	1010	0101	0101	0101	0101	1010	1010
_	1	0011	1100	1100	1100	_	0011	0011	1100	1100	0011	0011	8	911	118	1188
_	110	0110	1001	1001	1001		0110	0110	1001	1001	0110	0110	0110	0110	189	1001
_	000	1111	1111	0000	1111		0000	1111	1111	0000	000	1111	000	1111	1111	000
_	50	1010	1010	0101	1010		0101	1010	1010	0101	0101	1010	0101	1010	1010	0101
_	2	1100	1100	0011	1100		0011	1100	1100	811	811	1188	8	118	118	811
_	0110	<u>8</u>	1001	0110	1001		0110	1001	1001	0110	0110	1001	0110	100	100	0110
Ì																

HSLA

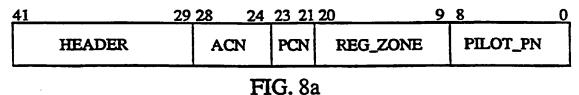
FIG. 6c

28

XMDZ- LOBX



ACCESS CHANNEL LONG CODE MASK



PUBLIC LONG CODE MASK

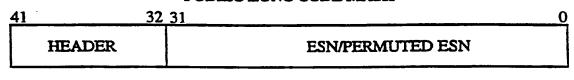


FIG. 8b

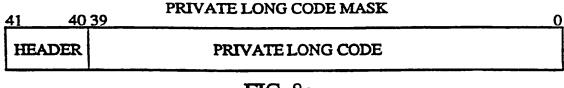
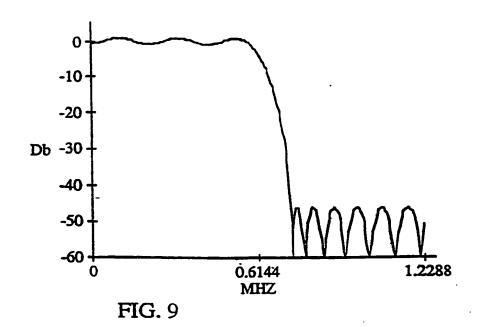


FIG. 8c







Europäisches Patentamt

European Patent Office

Office européen des brevets



EP 0 730 356 A3

(12)

EUROPEAN PATENT APPLICATION

(88) Date of publication A3: 20.11.1996 Bulletin 1996/47

(51) Int. Cl.⁶: **H04L 1/00**

(11)

(43) Date of publication A2: 04.09.1996 Bulletin 1996/36

(21) Application number: 96108491.0

(22) Date of filing: 19.01.1993

(84) Designated Contracting States:
AT BE CH DE DK ES FR GB GR IE IT LI LU MC NL
PT SE

(30) Priority: 16.01.1992 US 822164

(62) Application number of the earlier application in accordance with Art. 76 EPC: 93903541.6

(71) Applicant: QUALCOMM INCORPORATED San Diego, California 92121 (US)

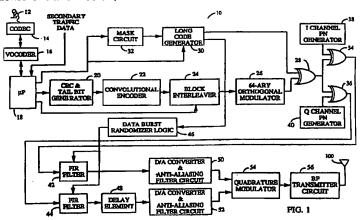
(72) Inventors:

- Padovani, Roberto San Diego, CA 92130 (US)
- Tiedemann, Edward, G., Jr.
 San Diego, CA 92122 (US)
- Weaver, Lindsay, A., Jr. Boulder, CO 80303 (US)
- Butler, Brian, K.
 Solana Beach, CA 92075 (US)
- (74) Representative: O'Sullivan, Gearoid Patrick et al Tomkins & Co.,
 5 Dartmouth Road Dublin 6 (IE)

(54) Communication system and method for transmitting data of different types

(57) The invention relates to a communication system in which transmission takes place according to a format which permits different types of data to be combined and transmitted within a single transmission. The novel feature is that the communication system transmits variable length frames of data in packets, and that a data combining and transmission sub-system (14, 16, 18, 20) is provided so that when a frame of data does not require a complete packet for transmission, the data

combining sub-system combines the frame of data with additional data to provide a complete packet. The data combining sub-system comprises input means for receiving the frame of data and the additional data and for combining the frame of data and the additional data to provide a complete packet responsive to a control signal, and control means for providing the control signal.





EUROPEAN SEARCH REPORT

Application Number EP 96 10 8491

Category	Citation of document with of relevant p	indication, where appropriate,	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.CL6)
X	EP-A-0 381 515 (NEC * page 2, column 2 column 3, line 21	line 15 - page 3,	1-4,9-11	H04L1/00
x	WO-A-88 04496 (PLES * page 2, line 11	SSEY OVERSEAS LIMITED) · line 13 *	1,9	
X	US-A-4 691 314 (BEI * column 5, line 63	GGINS ET AL.) B - column 6, line 17 *	1,9	
A	EP-A-0 189 695 (MOL	DLY MICHEL)	6,7,12,	
	* claim 1 *		13	
A	US-A-4 852 179 (FET * column 2, line 5	TE) - line 30 *	8	
				TECHNICAL FIELDS SEARCHED (Int.Cl.6)
				H04L H04J H04B
	The present search report has t	een drawn up for all claims	_	
	Place of search	Date of completion of the search		Excessioner
	THE HAGUE	13 September 199	96 Bis	chof, J-L
X : part Y : part doc: A : tech	CATEGORY OF CITED DOCUME icularly relevant if taken alone icularly relevant if combined with an ument of the same category inological background	E : earlier patent d	focument, but publi date I in the application for other reasons	ished on, or
O: non	-written disclosure	& : member of the		